## The Relational Model



Alan G. Labouseur, Ph.D.
Alan.Labouseur@Marist.edu

## The Relational Model

## Tables of rows and columns:

 Our AD\&D-like game

| CAP=\# select * CAP-\# from Players; |  |  |
| :---: | :---: | :---: |
| pid | name | rank |
| 1 | James |  Captain <br> $d$ Admiral |
| (2 rows) |  |  |
| CAP=\# select * |  |  |
| iid \| name | descr |  |  |
| A wand $\ldots$ <br> B gem $\ldots$ <br> C mace $\ldots$ <br> D sword $\ldots$. |  |  |
|  |  |  |
|  |  |  |
|  |  |  |
| (4 rows) |  |  |
| CAP=\# select * CAP-\# from Inventory; |  |  |
| pid \| iid | dateacquired |  |  |
| 1 A $2020-01-23$ <br> 1 B $\left\lvert\, \begin{array}{l}2020-01-23 \\ 2\end{array}\right.$ <br> 2 B $2020-01-23$ <br> 2 C $2020-01-23$ |  |  |
|  |  |  |
|  |  |  |
|  |  |  |
| (4 rows) |  |  |

## The Relational Model

## Tables of rows and columns: An e-commerce database ("CAP")



## The Relational Model

## Tables of rows and columns

People

| p |  | pre |  | firstName |  | lastName |  |  |  | homeCity |  | DOB |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 1 | Dr. | \| | Neil | \| | Peart |  | Ph.D. |  | Toronto |  | 1952-09-12 |
| 2 | । | Ms. | I | Regina | 1 | Schock | , |  | \| | Toronto |  | 1957-08-31 |
| 3 |  | Mr . | \| | Bruce | 1 | Crump | I | Jr. | 1 | Jacksonville |  | 1957-07-17 |
| 4 |  | Mr . | I | Todd | 1 | Sucherman |  |  |  | Chicago |  | 1969-05-02 |
| 5 | I | Mr . |  | Bernard | 1 | Purdie |  |  | 1 | Teaneck |  | 1939-06-11 |
| 6 | I | Ms. | I | Demetra | 1 | Plakas | I | Esq. | I | Santa Monica |  | 1960-11-09 |
| 7 | 1 | Ms. |  | Terri Lyne |  | Carrington | I |  |  | Boston |  | 1965-08-04 |
| 8 | 1 | Dr. |  | Bill | 1 | Bruford |  | Ph.D. |  | Kent |  | 1949-05-17 |
| 9 | 1 | Mr . |  | Alan | 1 | White | \| | III | 1 | Pelton |  | 1949-06-14 |

## The Relational Model

## Tables of rows and columns

People

|  |  | pres |  | firstName |  | lastName |  | suffix |  | homeCity |  | DOB |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 1 | Dr. | 1 | Neil |  | Peart |  | Ph.D. |  | Toronto |  | 1952-09-12 |
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## The Relational Model

## Tables of rows and columns

All entries in a single column are a single attribute and of a single data type.

| People |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| pid I prefix l firstName |  | I lastName | \| suffix | homeCity | DOB |
| 1 \| Dr. | I Neil | I Peart | I Ph.D. | \| Toronto | \| 1952-09-12 |
| 2 \| Ms. | \| Regina | I Schock | , | I Toronto | \| 1957-08-31 |
| 3 \| Mr. | I Bruce | I Crump | 1 Jr . | । Jacksonville | \| 1957-07-17 |
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| 9 \| Mr. | I Alan | \| White | 1 III | \| Pelton | \| 1949-06-14 |
|  |  |  |  |  |  |
| nteger Text |  |  |  |  |  |

## Relational Rules

## From Codd Himself.

These rules allow us to achieve excellence in database design across time and space.


## Relational Rules

## 1. The First Normal Form Rule

There can be no multi-valued attributes or values with internal structure at any intersection of a row and a column in a table.
In older terms: no repeating groups or no repeating fields.
Put another way, all values at the intersection or a row and a column must be atomic, meaning that they cannot be subdivided.

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People

| pid | name | profession | skills |
| :---: | :---: | :---: | :---: |
| 007 | Sean | spy | pronounces "S" like "Sh", charm |
| 008 | Roger | secret agent (on who's side?) | humour, stealth |
| 009 | Pierce | stiff-assed Brit | wit, hair |

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This is a violation of the 1 NF rule.

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In older terms: no repeating groups or no repeating fields.
Put another way, all values at the intersection or a row and a column must be atomic, meaning that they cannot be subdivided.

| pid | name | profession | skillı | skill2 |
| :---: | :---: | :---: | :---: | :---: |
| 007 | Sean | spy | pronounces "S" like "Sh" | charm |
| 008 | Roger | secret agent (on who's side?) | humour | stealth |
| 009 | Pierce | stiff-assed Brit | wit | hair |

A slight restructuring of table removes the 1 NF violation (but this is still bad design).

## Relational Rules

2. The Access Rows by Content Only Rule

We can only ask for ("query") data by what's there, never by where it is.

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We can only ask for ("query") data by what's there, never by where it is.

We can ask, "What is the name of pid 007?"
We cannot ask, "What is the name in the first row?"

| People |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
| pid | name | profession | skillı | skill2 |
| 007 | Sean | spy | pronounces "S" like "Sh" | charm |
| 008 | Roger | secret agent (on who's side?) | humour | stealth |
| 009 | Pierce | stiff-assed Brit | wit | hair |

Tables are sets. The elements of a set have no intrinsic order.
$\{\mathrm{a}, \mathrm{b}, \mathrm{c}\}=\{\mathrm{b}, \mathrm{a}, \mathrm{c}\}=\{\mathrm{c}, \mathrm{a}, \mathrm{b}\}$

## Relational Rules

## 2. The Access Rows by Content Only Rule

We can only ask for ("query") data by what's there, never by where it is.

We can ask, "What is the name of pid 7?" We cannot ask, "What is the name in ${ }^{\prime}$. Arst row?"


Tables are sets. The elements of a set have no intrinsic order. $\{\mathrm{a}, \mathrm{b}, \mathrm{c}\}=\{\mathrm{b}, \mathrm{a}, \mathrm{c}\}=\{\mathrm{c}, \mathrm{a}, \mathrm{b}\}$

## Relational Rules

3. The All Rows Must Be Unique Rule

Since tables are sets of rows and columns, and because the elements of a set have no intrinsic order, the only way we can insure our ability to get at every row in a table is for every row to be unique.

Were that not the case, some rows in the table would be indistinguishable. (Like crossing
 the streams, that would be bad.)

## Relational Rules

## 3. The All Rows Must Be Unique Rule

People


People


## Relational Rules

## Does this database obey the rules?



## Relational Rules

## Expanded Summary

- All entries in a table must be single-valued.
- Each column must have a distinct name.
- All values in a column are values of the same attribute.
- The order of columns is immaterial.
- Every row is distinct (unique).
- The order of rows is immaterial.

Keys

$$
\begin{array}{ll}
\text { Super Key } & \begin{array}{l}
\text { any field (column) or set of fields that } \\
\text { uniquely identify every row in a table }
\end{array}
\end{array}
$$

Candidate Key a minimal super key
Primary Key the chosen candidate key
Foreign Key a value in one table that must match the primary key of another table

## Super Key

## any field or set of fields that uniquely identify every row in a table



Products

| prodId । |  | name |  | city |  | OnHand |  | priceUSD |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| p01 |  | Heisenberg Compensator |  | Dallas |  | 47 |  | 67.50 |
| p02 | I | Universal Translator | I | Newark | I | 2399 |  | 5.50 |
| p03 | I | Commodore PET | I | Duluth | I | 1979 | I | 65.02 |
| p04 | I | LCARS module | I | Duluth | I | 3 | \| | 47.00 |
| p05 | \| | Remo drumhead | I | Dallas | 1 | 8675309 |  | 16.61 |
| p06 | I | Trapper Keeper | 1 | Dallas | I | 1982 |  | 2.00 |
| p07 |  | Flux Capacitor | I | Newark | I | 1007 |  | 1.00 |
| p08 | I | HAL 9000 memory core | 1 | Newark | 1 | 200 | I | 1.25 |
| p09 | I | Red Barchetta | I | Toronto | \| | 1 | \| | 379000.47 |

Agents

| pid | l paymentTerms | \| commissionPct |
| :---: | :--- | :--- |
| 2 | I Quarterly | \| |
| 3 | Annually | \| |
| 5 | I Monthly | 10.00 |
| 6 | \| Weekly | I |

## Orders



## Candidate Key a minimal super key



| pid | paymentTerms | commissionPct |
| :---: | :---: | :---: |
| 2 | Quarterly | 15.00 |
| 3 | Annually | 10.00 |
| 5 | Monthly | 2.00 |
| 6 | Weekly | 1.00 |

Orders

| orderNum | dateOrdered \| custId | agentId | prodId | quantityOrdered |  |  |  |  |  |  |  |  |  | totalUSD |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1011 | \| 2020-01-23 | I | 1 |  | 2 |  |  | I | 1100 |  | 58568.40 |
| 1012 | \| 2020-01-23 | I | 4 | I | 3 | I | p03 | I | 1200 | । | 74871.83 |
| 1015 | \| 2020-01-23 | I | 5 |  | 3 | I |  | I | 1000 | । | 15696.45 |
| 1016 | । 2020-01-23 | I | 8 | I | 3 | I | p01 | I | 1000 | । | 60750.00 |
| 1017 | \| 2020-02-14 | I | 1 | I |  | I |  | I | 500 |  | 25643.89 |
| 1018 | । 2020-02-14 | I | 1 | I |  | 1 |  | I | 600 |  | 22244.16 |
| 1019 | । 2020-02-14 | 1 | 1 | । |  | 1 | p02 | I | 400 | I | 1735.36 |
| 1020 | । 2020-02-14 | I | 4 | I |  | 1 |  | I | 600 | I | 575.76 |
| 1021 | \| 2020-02-14 | I | 4 | I | 5 | 1 | p01 | I | 1000 | I | 64773.00 |
| 1022 | । 2020-03-15 | I | 1 | , |  | 1 |  | I | 450 | I | 709.92 |
| 1023 | । 2020-03-15 | I | 1 | I |  | 1 |  | I | 500 | I | 6550.98 |
| 1024 | \| 2020-03-15 | I | 5 | I |  | 1 | p01 | I | 880 | I | 56133.00 |
| 1025 | I 2020-04-01 | I | 8 |  |  | 1 | p07 | I | 888 | । | 799.20 |
| 1026 | । 2020-05-01 | 1 | 8 |  | 5 | 1 | p03 | I | 808 | I | 47282.54 |

## Primary Key the chosen candidate key


Agents

| pid | paymentTerms | commissionPct |
| :---: | :--- | :--- |
| 2 | Quarterly | \| |
| 3 | Annually | \| |
| 5 | Monthly | 10.00 |
| 6 | \| Weekly | I |

## Products

| prodId | name | \| city |  | OnHand | priceUSD |
| :---: | :---: | :---: | :---: | :---: | :---: |
| p01 | \| Heisenberg Compensator | I Dallas | 1 | 47 | 167.50 |
| p02 | \| Universal Translator | I Newark | I | 2399 | 15.50 |
| p03 | I Commodore PET | I Duluth | 1 | 1979 | 65.02 |
| p04 | I LCARS module | I Duluth | I | 3 | 47.00 |
| p05 | I Remo drumhead | I Dallas | I | 8675309 | 16.61 |
| p06 | I Trapper Keeper | I Dallas | I | 1982 | 12.00 |
| p07 | I Flux Capacitor | I Newark | I | 1007 | 11.00 |
| p08 | । HAL 9000 memory core | I Newark | I | 200 | I 1.25 |
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## Keys

## Primary Key the chosen candidate key


Agents

| pid | paymentTerms | \| commissionPct |
| :---: | :--- | ---: |
| 2 | \| Quarterly | \| |


| Products |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| prodId | name | I city |  | tyOnHand | priceUSD |
| p01 | I Heisenberg Compensator | I Dallas | I | 47 | 67.50 |
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| p08 | I HAL 9000 memory core | 1 Newark | 1 | 200 | । 1.25 |
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## Foreign Key a value in one table that must match the primary key of another table



## Keys and Referential Integrity

The enforcement of the Primary Key (PK) - Foreign Key (FK) relationship is perhaps the most important aspect of Relational Databases. This property is called referential integrity. It insures consistency and accuracy, and thus leads to data quality. Data cannot become information without it.

Because of the importance of keys, it's critical that we - as database designers and data architects - never let end users control the content of key fields. For that reason, artificial keys are often a smart choice.

An artificial key is one that we make up. CWID is an example.

SELECT
FROM

## WHERE

## People



## SELECT some columns

## FROM

## WHERE

## People



## SELECT some columns <br> FROM some table

## WHERE



## SELECT some columns <br> FROM some table <br> WHERE some condition holds true

## People



## SELECT firstName <br> FROM People <br> wHERE homeCity = 'Toronto'

## People

| pid I pre | firstName | lastName | \| suffi | homeCity | DOB |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 1 \| Dr. | Neil | Peart | \| Ph.D. | \| Toronto | \| 1952-09-12 |
| 2 \| Ms. | Regina | Schock | I | I Toronto | \| 1957-08-31 |
| 3 \| Mr. | Bruce | Crump | 1 Jr. | \| Jacksonville | \| 1957-07-17 |
| 4 \| Mr. | I Todd | I Sucherman | I | \| Chicago | \| 1969-05-02 |
| 5 \| Mr. | \| Bernard | I Purdie | I | \| Teaneck | \| 1939-06-11 |
| 6 \| Ms. | Demetra | I Plakas | \\| Esq. | \| Santa Monica | \| 1960-11-09 |
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| 8 । Dr. | \| Bill | \| Bruford | I Ph.D. | I Kent | \| 1949-05-17 |
| 91 Mr . | \| Alan | I White | 1 III | I Pelton | \| 1949-06-14 |

## SQL

## SQL Script for AD\&D database: <br> Players, Items, and Inventory tables, and a few queries

```
create table Players (
    pid int not null,
    name text,
    rank text,
    primary key (pid)
);
insert into Players(pid, name, rank)
values (1, 'James', 'Captain'),
    (2, 'Leonard', 'Admiral');
select *
from Players;
create table Items (
    iid char(1) not null,
    name text,
    descr text,
    primary key (iid)
);
insert into Items (iid, name, descr)
values ('A', 'wand', '...''),
    ('B', 'gem', '...'),
    ('C', 'mace', '...'),
    ('D', 'sword', '...');
select *
from Items;
```

```
create table Inventory (
    pid int not null references Players(pid),
    iid char(1) not null references Items(iid),
    dateAcquired date,
    primary key(pid, iid)
);
insert into Inventory (pid, iid, dateAcquired)
values (1, 'A', '2020-01-23'),
            (1, 'В', '2020-01-23'),
            (2, 'В', '2020-01-23'),
            (2, 'C', '2020-01-23');
select *
from Inventory;
-- Players and their Items
select Players.name, Items.name
from Players inner join Inventory on Players.pid = Inventory.pid
            inner join Items on Inventory.iid = Items.iid;
-- Unused Items
select *
from Items
where iid not in (select iid
                                    from Inventory);
-- Item use count v1
select iid, count(iid)
from Inventory
group by iid
order by count(iid) DESC;
```

